



# GPU Teaching Kit

Accelerated Computing



## Module 10.4 – Parallel Computation Patterns (scan)

More on Parallel Scan

# Objective

- To learn more about parallel scan
  - Analysis of the work efficient kernel
  - Exclusive scan
  - Handling very large input vectors

# Work Analysis of the Work Efficient Kernel

- The work efficient kernel executes  $\log(n)$  parallel iterations in the reduction step
  - The iterations do  $n/2, n/4, \dots, 1$  adds
  - Total adds:  $(n-1) \rightarrow O(n)$  work
- It executes  $\log(n)-1$  parallel iterations in the post-reduction reverse step
  - The iterations do  $2-1, 4-1, \dots, n/2-1$  adds
  - Total adds:  $(n-2) - (\log(n)-1) \rightarrow O(n)$  work
- Both phases perform up to no more than  $2x(n-1)$  adds
- The total number of adds is no more than twice of that done in the efficient sequential algorithm
  - The benefit of parallelism can easily overcome the 2X work when there is sufficient hardware

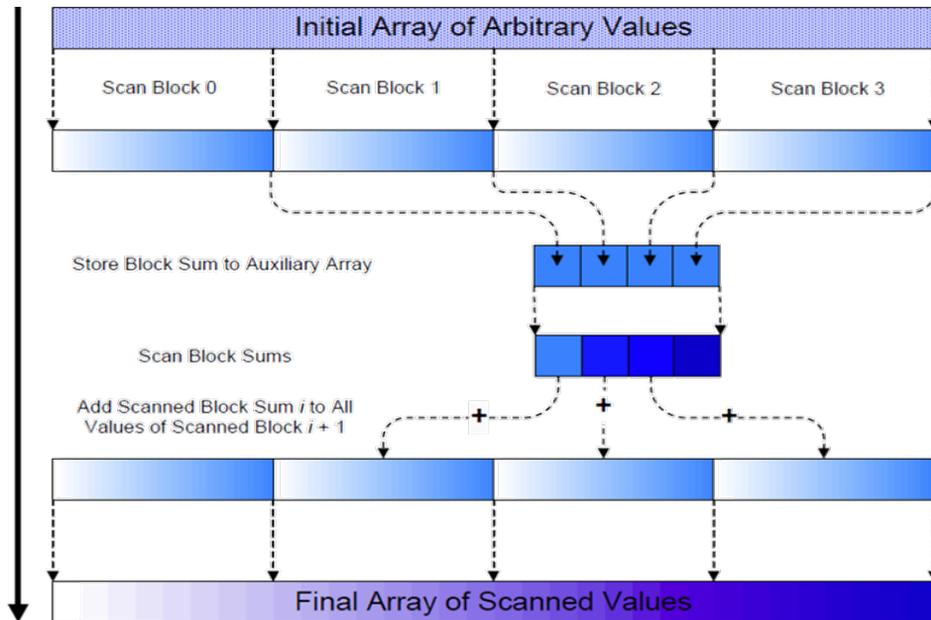
# Some Tradeoffs

- The work efficient scan kernel is normally more desirable
  - Better Energy efficiency
  - Less execution resource requirement
- However, the work inefficient kernel could be better for absolute performance due to its single-phase nature (forward phase only)
  - There is sufficient execution resource

# Handling Large Input Vectors

- Build on the work efficient scan kernel
- Have each section of  $2 \cdot \text{blockDim.x}$  elements assigned to a block
  - Perform parallel scan on each section
- Have each block write the sum of its section into a `Sum[]` array indexed by `blockIdx.x`
- Run the scan kernel on the `Sum[]` array
- Add the scanned `Sum[]` array values to all the elements of corresponding sections
- Adaptation of work inefficient kernel is similar.

# Overall Flow of Complete Scan



# Exclusive Scan Definition

**Definition:** The exclusive scan operation takes a binary associative operator  $\oplus$ , and an array of  $n$  elements

$$[x_0, x_1, \dots, x_{n-1}]$$

and returns the array

$$[0, x_0, (x_0 \oplus x_1), \dots, (x_0 \oplus x_1 \oplus \dots \oplus x_{n-2})].$$

**Example:** If  $\oplus$  is addition, then the exclusive scan operation on the array  $[3 \ 1 \ 7 \ 0 \ 4 \ 1 \ 6 \ 3]$ , would return  $[0 \ 3 \ 4 \ 11 \ 11 \ 15 \ 16 \ 22]$ .

# Why Use Exclusive Scan?

- To find the beginning address of allocated buffers
- Inclusive and exclusive scans can be easily derived from each other; it is a matter of convenience

[3 1 7 0 4 1 6 3]

Exclusive [0 3 4 11 11 15 16 22]

Inclusive [3 4 11 11 15 16 22 25]

# A Simple Exclusive Scan Kernel

- Adapt an inclusive, work inefficient scan kernel
- Block 0:
  - Thread 0 loads 0 into  $XY[0]$
  - Other threads load  $X[\text{threadIdx.x}-1]$  into  $XY[\text{threadIdx.x}]$
- All other blocks:
  - All thread load  $X[\text{blockIdx.x}*\text{blockDim.x}+\text{threadIdx.x}-1]$  into  $XY[\text{threadIdx.x}]$
- Similar adaption for work efficient scan kernel but ensure that each thread loads two elements
  - Only one zero should be loaded
  - All elements should be shifted to the right by only one position

Read the Harris article (Parallel Prefix Sum with CUDA) for a more intellectually interesting approach to exclusive scan kernel implementation.



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